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# A RECURSIVE ALGORITHM TO COMPUTE A BASIS OF A SIMILARITY

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## Abstract

This paper studies some theory and methods to build a representation theorem basis of a similarity from the basis of its subsimilarities, providing an alternative recursive method to compute the basis of a similarity.

**Keywords:** T-indistinguishabilities representation theorem, similarity, basis of a similarity, dimension of a similarity, similarity representation theorem.

## 1 INTRODUCTION

Similarity fuzzy relations were introduced by Zadeh [1971] [10] to represent a degree of equality or closeness between the elements of a universe. They generalises the equivalence relations; in fact, the similarities are the only T-indistinguishabilities that satisfy that all their alpha cuts are crisp equivalence relations. Similarities are not only a powerful tool to represent equality information, but they are also useful to classify the universe into clusters with uncertainty.

The Valverde's representation theorem of T-indistinguishabilities is one of the strongest theorems in fuzzy logic theory. It opened an interesting area of investigation on T-indistinguishabilities, the computation of its basis, specially for both the minimum t-norm and Archimedean t-norms, and the study of special T-indistinguishabilities such as the one dimensional ones.

On the other hand, the computational investigation on the computation of T-transitive closures, or the search of the structure of a similarity, specially its decomposition into subsimilarities, were not related in previous studies with the computation of the basis of a similarity.

Joan Jacas [4] studied two algorithms to compute basis of a similarity in 1990, but did not consider the decomposition of similarities to do it. The aim of this paper is to introduce a new approach to compute them, providing useful theory and methods for both better understanding the concept of structure of a similarity and the Valverde's representation theorem for similarities and the computation of their basis using a new decomposition approach.

## 2 PRELIMINARIES

Let  $X = \{x_1, \dots, x_n\}$  be a finite universe.

Let  $T$  be a t-norm. A  $T$ -indistinguishability operator  $E$  on  $X$  is a fuzzy relation  $E: X \times X \rightarrow [0, 1]$ , satisfying for all  $x, y, z$  in  $X$ :

1.  $E(x, x) = 1$  (Reflexivity)
2.  $E(x, y) = E(y, x)$  (Symmetry)
3.  $T(E(x, y), E(y, z)) \leq E(x, z)$  (T-transitivity)

**Definition 2.2.** A **similarity** [Zadeh, 1971] [10] is a reflexive, symmetric and min-transitive fuzzy relation, it is, a similarity is a Min-indistinguishability operator.

**Notation 2.1.**

We can denote  $x_{ij} = E(x_i, x_j)$ .

**Lemma 2.1.** Let  $\pi$  be a permutation on  $X$ . If  $E$  is a similarity on  $X$ , then the fuzzy relation  $P_\pi(E)$  is also a fuzzy similarity.

**Proof.** It is obvious.  $P_\pi(E)$  is reflexive and symmetric. If  $x_{ij} \geq \min\{x_{ik}, x_{kj}\}$  for all  $i, j, k$  then  $x_{rs} = x_{\pi(i)\pi(j)} \geq \min\{x_{\pi(i)\pi(k)}, x_{\pi(k)\pi(j)}\} = T\{x_{rk}, x_{ks}\}$  for all  $1 \leq r, s, k \leq n$ .

### 2.1. CONSTRUCTION OF A FUZZY SIMILARITY FROM SUBSIMILARITIES

Let  $C$  and  $D$  be two similarities on two disjoint sets  $X_1$  and  $X_2$ . A similarity relation  $E(F; C, D)$  on  $X_1 \cup X_2$  can be built with the following shape:

$$E(F; C, D) = \begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix}$$

A method for giving the bridging values  $f_{ij}$  in  $F$ , (when  $j \leq \text{card}(X_1) < i$ ) is the assignation of a unique value  $f$ , in all the  $\text{card}(X_1) \times \text{card}(X_2)$  values in  $F$ . This value must be chosen in an interval  $[0, a]$  where  $a = \min\{\min(C), \min(D)\}$ . The values in  $F^T$  are the symmetric values  $f$  of the computed  $F$ .

So the computed values in  $F$  are equal and satisfy that  $f \leq \min\{\min(C), \min(D)\}$ .

**Lemma 2.1.1.** [5] Lee [2001]

If  $C$  and  $D$  are fuzzy similarities, then  $E(f; C, D)$  is also a fuzzy similarity,  $\forall f \in [0, \min(\min(C), \min(D))]$ .

**Lemma 2.1.2.** [5] Lee [2001]

If  $S_{n \times n}$  is a fuzzy similarity, then there exists a decomposition such that  $S_{n \times n} = P_{\pi}(E(f; C_{n1 \times n1}, D_{n2 \times n2}))$ .

## 2.2. THE REPRESENTATION THEOREM OF T-INDISTINGUISHABILITY OPERATORS

The representation theorem allow us to generate a  $T$ -indistinguishability operator on a set  $X$  from a family of subsets on  $X$ , and reciprocally states that every  $T$ -indistinguishability can be obtained in this form.

**Definition 2.2.1** The residuation  $\vec{T}$  or quasi inverse of a t-norm  $T$  is the map  $\vec{T}: [0, 1] \times [0, 1] \rightarrow [0, 1]$  defined for all  $x, y$  in  $[0, 1]$  by

$$\vec{T}(x | y) = \sup\{\alpha \in [0, 1] | T(x, \alpha) \leq y\}.$$

The residuation  $\vec{T}$  of a t-norm  $T$  is a  $T$ -preorder (reflexive and  $T$ -transitive) on  $[0, 1]$ , and then it is a useful operator to generate implication relations from fuzzy sets on  $X$ .

Note that  $\overrightarrow{Min}(x | y) = \begin{cases} 1 & \text{if } x \leq y \\ y & \text{if } x > y \end{cases}$ .

**Definition 2.2.2** The biresiduation  $\vec{T}$  (or also  $E_T$ ) of a t-norm  $T$  is the map  $\vec{T}: [0, 1] \times [0, 1] \rightarrow [0, 1]$  defined for all  $x, y$  in  $[0, 1]$  by

$$\begin{aligned} \vec{T}(x, y) &= E_T(x, y) = \min(\vec{T}(x | y), \vec{T}(y | x)). \\ \text{Note that } \overrightarrow{Min}(x, y) &= \overrightarrow{Min}(\max(x, y), \min(x, y)) \\ &= \begin{cases} 1 & \text{if } x = y \\ \min(x, y) & \text{if } x \neq y \end{cases} \end{aligned}$$

The biresiduation  $\vec{T}$  of a t-norm  $T$  is a  $T$ -indistinguishability operator on  $[0, 1]$ , and then it is a useful operator to generate  $T$ -indistinguishability relations on  $X$  from two fuzzy sets on  $X$ .

**Lemma 2.2.1.**

Let  $\mu$  be a fuzzy set on  $X$ , and  $T$  a continuous t-norm. The fuzzy relation  $E_{\mu}$  on  $X$  defined for all  $x, y$  in  $X$  by  $E_{\mu}(x, y) = E_T(\mu(x), \mu(y))$  is a  $T$ -indistinguishability operator.

Note that  $E_{\mu}$  is a one-dimensional  $T$ -indistinguishability generated by a basis of one fuzzy set  $\mu$ .

**Lemma 2.2.2.**

Let  $(E_i)_{i \in I}$  be a family of  $T$ -indistinguishability operators on a set  $X$ . The relation  $E$  on  $X$  defined for all  $x, y$  in  $X$  by

$$E(x, y) = \inf_{i \in I} E_{\mu_i}(x, y)$$

is a  $T$ -indistinguishability operator.

The next theorem is crucial to understand the structure of a  $T$ -indistinguishability operator. It allows us to generate

$T$ -indistinguishabilities from a family of fuzzy sets, and reciprocally that any  $T$ -indistinguishability can be generated from a family of fuzzy sets.

**Representation theorem 2.2.1.** [9]

Let  $R$  be a fuzzy relation on  $X$  and  $T$  a continuous t-norm.  $R$  is a  $T$ -indistinguishability operator if and only if there exists a family  $(\mu_i)_{i \in I}$  of fuzzy sets on  $X$  such that for all  $x, y$  in  $X$

$$R(x, y) = \inf_{i \in I} E_{\mu_i}(x, y).$$

**Definition 2.2.1. Dimension and basis of a  $T$ -indistinguishability operator**

The dimension of a  $T$ -indistinguishability operator  $E$  is the minimal of the cardinalities  $d$  of the generating families of  $E$ , in the sense of the representation theorem.

That minimal basis of generating fuzzy sets is called a basis of the  $T$ -indistinguishability operator.

## 3 DECOMPOSITION OF SIMILARITIES AND REPRESENTATION THEOREM

This chapter gives some theoretical results toward a method to compute a basis of a similarity from the bases of its subsimilarities.

**Lemma 3.1.** Let  $S$  be the following similarity on  $X_1 \cup X_2$ ,

$$S(F; C, D) = \begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix}$$

Then  $\dim(S) \geq \dim(C)$ .

**Lemma 3.2.**

Let  $(\mu_i)_{i \in I}$  be a representation theorem basis of a similarity  $C$  of dimension  $r$  on a finite set  $X_1$ . Let  $(\gamma_i)_{i \in J}$  be a representation theorem basis of a similarity  $D$  of dimension  $s$  on a finite set  $X_2$ .

Let  $S$  be the following similarity on  $X_1 \cup X_2$ ,  $S(F; C, D) =$

$$\begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix} \text{ where } F = \begin{pmatrix} f & \cdots & f \\ \vdots & \ddots & \vdots \\ f & \cdots & f \end{pmatrix} \text{ is a } \text{card}(X_1) \times$$

$\text{card}(X_2)$  matrix.

Suppose that  $r > s$  and  $f < \min(C)$  and  $f \leq \min(D)$  then  $\dim(S) \leq \dim(C)$  and a generator set of  $S$  is

$$\left\{ \begin{pmatrix} \mu_1 \\ \gamma_1 \end{pmatrix}, \dots, \begin{pmatrix} \mu_s \\ \gamma_s \end{pmatrix}, \begin{pmatrix} \mu_{s+1} \\ F' \end{pmatrix}, \dots, \begin{pmatrix} \mu_r \\ F' \end{pmatrix} \right\} \text{ where } F' = \begin{pmatrix} f \\ \vdots \\ f \end{pmatrix} \text{ is a } \text{card}(X_2) \times 1 \text{ matrix.}$$

**Proposition 3.1.**

Let  $(\mu_i)_{i \in I}$  be a representation theorem basis of a similarity C of dimension  $r$  on a finite set  $X_1$ . Let  $(\gamma_i)_{i \in J}$  be a representation theorem basis of a similarity D of dimension  $s$  on a finite set  $X_2$ .

Let S be the following similarity on  $X_1 \cup X_2$ ,  $S(F; C, D) = \begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix}$  where  $F = \begin{pmatrix} f & \cdots & f \\ \vdots & \ddots & \vdots \\ f & \cdots & f \end{pmatrix}$  is a  $\text{card}(X_1) \times \text{card}(X_2)$  matrix.

Suppose that  $r > s$  and  $f < \min(C)$  and  $f \leq \min(D)$  then

- 1)  $\dim(S) = \dim(C) = r$
- 2) a basis of S is  $\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F' \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ F' \end{pmatrix} \right\}$

where  $F' = \begin{pmatrix} f \\ \vdots \\ f \end{pmatrix}$  is a  $\text{card}(X_2) \times 1$  matrix.

**Proof:**

By Lemma 3.1  $\dim(S) \geq \dim(C)$ .

By Lemma 3.2  $\dim(S) \leq \dim(C)$ . Also  $\dim(S) = \dim(C) = r$  and

$\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F' \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ F' \end{pmatrix} \right\}$  is a basis of  $S(F; C, D)$ .  $\square$

**Corollary 3.1.**

Let  $(\mu_i)_{i \in I}$  be a representations theorem basis of a similarity C on a finite set  $X = \{x_1, \dots, x_n\}$ .

Let S be the similarity on  $X \cup \{x_{n+1}\}$  such that  $S(F; C, 1) = \begin{pmatrix} C & F'^T \\ F' & 1 \end{pmatrix}$  where  $F' = (f \dots f)$  is a  $1 \times \text{card}(X)$  matrix.

If  $f < \min(C)$ , then a basis of S is

$$\left\{ \begin{pmatrix} \mu_1 \\ \vdots \\ \mu_1(x_n) \\ f \end{pmatrix}, \dots, \begin{pmatrix} \mu_{\dim(C)} \\ \vdots \\ \mu_{\dim(C)}(x_n) \\ f \end{pmatrix} \right\} = \left\{ \begin{pmatrix} \mu_1(x_1) \\ \vdots \\ \mu_1(x_n) \\ f \end{pmatrix}, \dots, \begin{pmatrix} \mu_{\dim(C)}(x_1) \\ \vdots \\ \mu_{\dim(C)}(x_n) \\ f \end{pmatrix} \right\}$$

**Example 3.1.**

Let S be the similarity  $\begin{pmatrix} 1 & a & b & c \\ a & 1 & b & c \\ b & b & 1 & c \\ c & c & c & 1 \end{pmatrix}$ , already ordered

with  $c < b < a$ .

$$S = \begin{pmatrix} 1 & a & b & c \\ a & 1 & b & c \\ b & b & 1 & c \\ c & c & c & 1 \end{pmatrix},$$

and a base of  $\begin{pmatrix} 1 & a & b \\ a & 1 & b \\ b & b & 1 \end{pmatrix}$  is  $\left\{ \begin{pmatrix} 1 \\ a \\ b \end{pmatrix} \right\}$ , then by the theorem 3.1 a basis of S is

$$\left\{ \begin{pmatrix} 1 \\ a \\ b \end{pmatrix} \right\} \times \{c\} = \left\{ \begin{pmatrix} 1 \\ a \\ b \\ c \end{pmatrix} \right\}.$$

Note that S is one dimensional.

**Example 3.2.**

Let S be the similarity  $\begin{pmatrix} 1 & a & a & b \\ a & 1 & a & b \\ a & a & 1 & b \\ b & b & b & 1 \end{pmatrix}$ , already ordered

with  $b < a$ .

$$S = \begin{pmatrix} 1 & a & a & b \\ a & 1 & a & b \\ a & a & 1 & b \\ b & b & b & 1 \end{pmatrix},$$

A basis of  $\begin{pmatrix} 1 & a & a \\ a & 1 & a \\ a & a & 1 \end{pmatrix}$  is  $\left\{ \begin{pmatrix} 1 \\ a \\ a \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \\ a \end{pmatrix} \right\}$

Then, as  $b < \min \begin{pmatrix} 1 & a & a \\ a & 1 & a \\ a & a & 1 \end{pmatrix}$ , by theorem 3.1, a basis of S

$$\text{is } \left\{ \begin{pmatrix} 1 \\ a \\ a \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \\ a \end{pmatrix} \right\} \times \{b\} = \left\{ \begin{pmatrix} 1 \\ a \\ a \\ b \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \\ a \\ b \end{pmatrix} \right\}$$

**Lemma 3.2.**

Let  $(\mu_i)_{i \in I}$  be a representation theorem basis of a similarity C of dimension  $r$  on a finite set  $X_1$ . Let  $(\gamma_i)_{i \in J}$  be a representation theorem basis of a similarity D of dimension  $s$  on a finite set  $X_2$ .

Let S be the following similarity on  $X_1 \cup X_2$ ,  $S(F; C, D) = \begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix}$ , where  $F = \begin{pmatrix} f & \cdots & f \\ \vdots & \ddots & \vdots \\ f & \cdots & f \end{pmatrix}$  is a  $\text{card}(X_1) \times \text{card}(X_2)$  matrix

Suppose that  $r > s, f = \min(C)$  and  $f \leq \min(D)$ , then:

- 1)  $\dim(S) > \dim(C) = r$ .

**Lemma 3.3.**

Let  $(\mu_i)_{i \in I}$  be a representation theorem basis of a similarity C of dimension  $r$  on a finite set  $X_1$ . Let  $(\gamma_i)_{i \in J}$  be a representation theorem basis of a similarity D of dimension  $s$  on a finite set  $X_2$ .

Let  $S$  be the following similarity on  $X_1 \cup X_2$ ,  $S(F; C, D) = \begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix}$ ,

where  $F = \begin{pmatrix} f & \dots & f \\ \vdots & \ddots & \vdots \\ f & \dots & f \end{pmatrix}$  is a  $\text{card}(X_1) \times \text{card}(X_2)$  matrix

Suppose that  $r > s$ ,  $f = \min(C)$  and  $f \leq \min(D)$ , then:

- 1)  $\dim(S) \leq \dim(C) + 1 = r + 1$
- 2) A generator set of  $S$  is A basis of  $S$  is  $\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F' \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ F' \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\}$ ,

where  $(1) = \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}$  is a  $\text{card}(X_1) \times 1$  matrix and  $F' = \begin{pmatrix} f \\ \vdots \\ f \end{pmatrix}$  is a  $\text{card}(X_2) \times 1$  matrix.

**Proof:**

Now  $\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F' \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ F' \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\}$  is a generator set of  $S$ , because

$$\inf \left\{ \inf_{1 \leq i \leq s} E_{\begin{pmatrix} (\mu_i) \\ (\gamma_i) \end{pmatrix}}, \inf_{r+1 \leq j \leq r} E_{\begin{pmatrix} (\mu_j) \\ F' \end{pmatrix}} \right\} = \begin{pmatrix} C & F^* \\ F^* & D \end{pmatrix} = S(F; C, D).$$

Also  $\begin{pmatrix} (1) \\ F' \end{pmatrix} o_{\overline{\min}} ((1), F') = \begin{pmatrix} (1) & F^T \\ F & (1) \end{pmatrix}$ , so the infimum of all the similarities is  $\begin{pmatrix} C & F^T \\ F & D \end{pmatrix} = S(F; C, D)$ .  $\square$

### Proposition 3.2.

Let  $(\mu_i)_{i \in I}$  be a representation theorem basis of a similarity  $C$  of dimension  $r$  on a finite set  $X_1$ . Let  $(\gamma_i)_{i \in J}$  be a representation theorem basis of a similarity  $D$  of dimension  $s$  on a finite set  $X_2$ .

Let  $S$  be the following similarity on  $X_1 \cup X_2$ ,  $S(F; C, D) = \begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix}$ , where  $F = \begin{pmatrix} f & \dots & f \\ \vdots & \ddots & \vdots \\ f & \dots & f \end{pmatrix}$  is a  $\text{card}(X_1) \times \text{card}(X_2)$  matrix

Suppose that  $r > s$ ,  $f = \min(C)$  and  $f \leq \min(D)$ , then:

- 1)  $\dim(S) = \dim(C) + 1 = r + 1$
- 2) A basis of  $S$  is:

$$\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F' \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ F' \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\}$$

where  $(1) = \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}$  is a  $\text{card}(X_1) \times 1$  matrix and  $F' = \begin{pmatrix} f \\ \vdots \\ f \end{pmatrix}$  is a  $\text{card}(X_2) \times 1$  matrix.

**Proof:**

By lemma 3.2  $\dim(S) > \dim(C) = r$ , and by lemma 3.3  $\dim(S) \leq \dim(C) + 1 = r + 1$ , then  $\dim(S) = \dim(C) + 1 = r + 1$ , and a basis of  $S$  is  $\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F' \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ F' \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\}$ ,

### Example 3.5

Let  $S$  be the similarity  $\begin{pmatrix} 1 & a & b & b \\ a & 1 & b & b \\ b & b & 1 & b \\ b & b & b & 1 \end{pmatrix}$ , already ordered with  $b < a$ .

$S = \begin{pmatrix} \begin{pmatrix} 1 & a & b \\ a & 1 & b \\ b & b & 1 \end{pmatrix} & \begin{pmatrix} b \\ b \\ b \end{pmatrix} \\ \begin{pmatrix} b \\ b \\ b \end{pmatrix} & 1 \end{pmatrix}$ , and a basis of  $\begin{pmatrix} 1 & a & b \\ a & 1 & b \\ b & b & 1 \end{pmatrix}$ , is  $\left\{ \begin{pmatrix} 1 \\ a \\ b \end{pmatrix} \right\}$ , then by the propositions 3.1 and 3.2 a generator set of  $S$  is  $= \left\{ \begin{pmatrix} (1) \\ a \\ b \\ b \end{pmatrix}, \begin{pmatrix} (1) \\ 1 \\ 1 \\ b \end{pmatrix} \right\}$ .

### Proposition 3.3.

Let  $(\mu_i)_{i \in I}$  be a representation theorem basis of a similarity  $C$  of dimension  $r$  on a finite set  $X_1$ . Let  $(\gamma_i)_{i \in J}$  be a representation theorem basis of a similarity  $D$  of dimension  $s$  on a finite set  $X_2$ .

Let  $S$  be the following similarity on  $X_1 \cup X_2$ ,  $S(F; C, D) = \begin{pmatrix} \boxed{C} & \boxed{F^T} \\ \boxed{F} & \boxed{D} \end{pmatrix}$ ,

where  $F = \begin{pmatrix} f & \dots & f \\ \vdots & \ddots & \vdots \\ f & \dots & f \end{pmatrix}$  is a  $\text{card}(X_1) \times \text{card}(X_2)$  matrix.

Suppose that  $r = s$  ( $\dim(C) = \dim(D)$ ) and  $f < \min(C, D)$ , then a generator set of  $S$  is:

$$\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\}$$

and  $\dim(S) \leq r + 1$ .

**Proof:**

Let  $\{\mu_1, \dots, \mu_r\}$  be a representation theorem basis of a similarity  $C$ , so

$$C = \inf_{i \in I} E_{\mu_i} \text{ where } I = \{1, \dots, r\}.$$



In the same way, let  $\{\gamma_1, \dots, \gamma_s\}$  be a representation theorem basis of a similarity D, so

$$D = \inf_{j \in J} E_{\gamma_j} \text{ where } J = \{1, \dots, r\}$$

Now, for all  $1 \leq i \leq r$ ,

$$\begin{pmatrix} (\mu_i) \\ (\gamma_i) \end{pmatrix} o_{\overline{Min}}((\mu_i) \quad (\gamma_i)) = \begin{pmatrix} (E_{\mu_i}) & (x_{nm}) \\ (x_{mn}) & (E_{\gamma_i}) \end{pmatrix}$$

$$\text{Also } \begin{pmatrix} (1) \\ F' \end{pmatrix} o_{\overline{Min}}((1) \quad F') = \begin{pmatrix} (1) & F^T \\ F & (1)^T \end{pmatrix},$$

So the infimum of all the similarities is  $\begin{pmatrix} C & F^T \\ F & D \end{pmatrix} = S(F; C, D)$ , and then  $\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\}$  is a generator set of S.  $\square$

#### Example 3.4.

Let S be the similarity  $\begin{pmatrix} 1 & a & b & b \\ a & 1 & b & b \\ b & b & 1 & a \\ b & b & a & 1 \end{pmatrix}$ , already ordered,

with  $b < a$ .

$$S = \begin{pmatrix} \begin{pmatrix} 1 & a \\ a & 1 \end{pmatrix} & \begin{pmatrix} b & b \\ b & b \end{pmatrix} \\ \begin{pmatrix} b & b \\ b & b \end{pmatrix} & \begin{pmatrix} 1 & a \\ a & 1 \end{pmatrix} \end{pmatrix}$$

A basis of  $\begin{pmatrix} 1 & a \\ a & 1 \end{pmatrix}$  is  $\left\{ \begin{pmatrix} 1 \\ a \end{pmatrix} \right\}$ , then by the proposition 3.3 a generator set (and a basis in this example) of S is

$$\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\} = \left\{ \begin{pmatrix} 1 \\ a \end{pmatrix}, \begin{pmatrix} 1 \\ b \end{pmatrix} \right\}.$$

Note that there exists other bases of S, for example

$$\left\{ \begin{pmatrix} 1 \\ a \end{pmatrix}, \begin{pmatrix} b \\ a \end{pmatrix} \right\}.$$

Let S be the similarity  $\begin{pmatrix} 1 & a & b & b \\ a & 1 & b & b \\ b & b & 1 & a \\ b & b & a & 1 \end{pmatrix}$ , already ordered,

with  $b < a$ .

$$S = \begin{pmatrix} \begin{pmatrix} 1 & a \\ a & 1 \end{pmatrix} & \begin{pmatrix} b & b \\ b & b \end{pmatrix} \\ \begin{pmatrix} b & b \\ b & b \end{pmatrix} & \begin{pmatrix} 1 & a \\ a & 1 \end{pmatrix} \end{pmatrix}$$

A base of  $\begin{pmatrix} 1 & a \\ a & 1 \end{pmatrix}$  is  $\left\{ \begin{pmatrix} 1 \\ a \end{pmatrix} \right\}$ ,

Then by proposition 3.3 a basis of S is

$$= \left\{ \begin{pmatrix} 1 \\ a \end{pmatrix}, \begin{pmatrix} b \\ a \end{pmatrix} \right\}.$$

## 4 GENERATING A DECOMPOSITION OF A GIVEN SIMILARITY

Input: A similarity S on X.

Output:

- Two similarities  $S_1$ , and  $S_2$  on  $X_1$  and  $X_2$  such that  $X = X_1 \cup X_2$
- A bridging value  $f$  such that  $f < \min(S_1)$  and  $f \leq \min(S_2)$

Algorithm 5.1:

1. Sort the universe X having descending columns under the diagonal using algorithm 4.1
2. Compute the frequency  $k$  of the lowest value  $v$  in the first column
3. If  $k \geq \text{card}(X) - 1$  then STOP.  $S_1 = S$ .
4. Let  $X_1$  be  $\{x_1, \dots, x_{k+1}\}$  and let  $X_2$  be  $\{x_{k+2}, \dots, x_{\text{card}(X)}\}$
5. Let  $S_1 = S|_{X_1}$ ,  $S_2 = S|_{X_2}$ , and  $f = v$ .

## 5 A RECURSIVE ALGORITHM TO COMPUTE A GENERATING SET OF A SIMILARITY

Input: a similarity S on a universe X.

Output: A generating set of S, and so, an upper bound of the dimension of S.

A recursive algorithm:

- 1) If  $\text{card}(X) \leq 2$  then STOP.  $\dim(S) = 1$  and a basis of S is the first column.
- 2) Decompose S into  $S_1$ , and  $S_2$  using algorithm 5.1. Note that  $f < \min(S_1)$
- 3) IF  $\text{card}(X_2) = 1$  then use corollary 3.1:
  - a. Recursively compute a generating set of  $S_1, (\mu_i)_{i \in I}$
  - b. A generating set of S is  $\{(\mu_i)_{i \in I}\} \times \{f\}$
  - c. STOP
- 4) Recursively compute a generating set of  $S_1, (\mu_i)_{i \in I}$  and  $S_2, (\gamma_i)_{i \in J}$
- 5) IF  $\dim(S_1) > \dim(S_2)$  -or vice versa- then use proposition 3.1
  - a. A generating set of S is  $\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F \end{pmatrix}, \dots \right\}$
  - b. STOP
- 6) IF  $\dim(S_1) = \dim(S_2)$  then use proposition 3.3
  - a. A generating set of S is  $\left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (1) \\ F \end{pmatrix} \right\}$
  - b. STOP

**Example 6.1.**

Let  $S$  be the sorted similarity with descending columns under the diagonal given by the following matrix:

$$S = \begin{pmatrix} 1 & 0,9 & 0,4 & 0,4 & 0,3 \\ 0,9 & 1 & 0,4 & 0,4 & 0,3 \\ 0,4 & 0,4 & 1 & 0,7 & 0,3 \\ 0,4 & 0,4 & 0,7 & 1 & 0,3 \\ 0,3 & 0,3 & 0,3 & 0,3 & 1 \end{pmatrix}$$

It will be computed the algorithm 6.1 step by step.

A decomposition of  $S$  is given by algorithm 5.1 is

$$S = \begin{pmatrix} 1 & 0,9 & 0,4 & 0,4 & 0,3 \\ 0,9 & 1 & 0,4 & 0,4 & 0,3 \\ 0,4 & 0,4 & 1 & 0,7 & 0,3 \\ 0,4 & 0,4 & 0,7 & 1 & 0,3 \\ 0,3 & 0,3 & 0,3 & 0,3 & 1 \end{pmatrix}, S_1 = \begin{pmatrix} 1 & 0,9 & 0,4 & 0,4 \\ 0,9 & 1 & 0,4 & 0,4 \\ 0,4 & 0,4 & 1 & 0,7 \\ 0,4 & 0,4 & 0,7 & 1 \end{pmatrix}$$

Now that  $S$  is ordered and decomposed, we can build its basis as follows:

A basis of  $\begin{pmatrix} 1 & 0,9 \\ 0,9 & 1 \end{pmatrix}$  is  $\left\{ \begin{pmatrix} 1 \\ 0,9 \end{pmatrix} \right\}$ . A basis of  $\begin{pmatrix} 1 & 0,7 \\ 0,7 & 1 \end{pmatrix}$  is  $\left\{ \begin{pmatrix} 1 \\ 0,7 \end{pmatrix} \right\}$ , and  $\dim C = 2 = \dim D$ ,  $f = 0,4 < \min(C, D)$ .

Then by proposition 3.3, a generator set of  $S_1 =$

$$\left\{ \begin{pmatrix} 1 & 0,9 \\ 0,9 & 1 \end{pmatrix}, \begin{pmatrix} 0,4 & 0,4 \\ 0,4 & 0,4 \end{pmatrix}, \begin{pmatrix} 1 & 0,7 \\ 0,7 & 1 \end{pmatrix} \right\}$$

$$\text{is } \left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (1) \\ F' \end{pmatrix} \right\} = \left\{ \begin{pmatrix} 1 \\ 0,9 \\ 1 \\ 0,7 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \\ 0,4 \\ 0,4 \end{pmatrix} \right\}$$

And then, by proposition 3.1, a generator set of  $S =$

$$\begin{pmatrix} S_1 & 0,3 \\ 0,3 & (1) \end{pmatrix} \text{ is } \left\{ \begin{pmatrix} (\mu_1) \\ (\gamma_1) \end{pmatrix}, \dots, \begin{pmatrix} (\mu_s) \\ (\gamma_s) \end{pmatrix}, \begin{pmatrix} (\mu_{s+1}) \\ F' \end{pmatrix}, \dots, \begin{pmatrix} (\mu_r) \\ F' \end{pmatrix} \right\} = \left\{ \begin{pmatrix} 1 \\ 0,9 \\ 1 \\ 0,7 \\ 0,3 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \\ 0,4 \\ 0,4 \\ 0,3 \end{pmatrix} \right\}$$

**6 CONCLUDING REMARKS**

This paper's main contribution is a method to compute a representation theorem basis of a similarity from the bases of its subsimilarities.

These results can be used to propose an alternative algorithm to build a basis of similarities.

The bases of all structures of similarities with dimension four are computed in the examples using the new construction method.

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